

Definitions related to regular expressions
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Definition 1 Let A be an alphabet. A regular expression over A is defined recursively as follows:

1. \emptyset is a regular expression over A .
2. Λ is a regular expression over A .
3. $\forall a \in A$, a is a regular expression over A .
4. If r is a regular expression over A , then (r) is a regular expression over A .
5. If r is a regular expression over A , then r^* is a regular expression over A .
6. If r_1 and r_2 are regular expressions over A , then $r_1 + r_2$ is a regular expression over A .
7. If r_1 and r_2 are regular expressions over A , then $r_1 r_2$ is a regular expression over A .

Definition 2 For a regular expression r over alphabet A , we define the language defined by r , $\text{lang}(r)$, recursively as follows.

1. $\text{lang}(\emptyset) = \{\}$.
2. $\text{lang}(\Lambda) = \{\Lambda\}$.
3. If $a \in A$, then $\text{lang}(a) = \{a\}$.
4. If r is a regular expression over A , then $\text{lang}((r)) = \text{lang}(r)$.
5. If r is regular expression over A , then $\text{lang}(r^*) = (\text{lang}(r))^*$.
6. If r_1 and r_2 are regular expressions over A , then

$$\text{lang}(r_1 + r_2) = \text{lang}(r_1) \cup \text{lang}(r_2).$$

7. If r_1 and r_2 are regular expressions over A , then

$$\text{lang}(r_1 r_2) = \text{lang}(r_1) \text{lang}(r_2).$$

Definition 3 Let L be a language over alphabet A . L is regular if there exists a regular expression r such that $\text{lang}(r) = L$.

Note. There is an order of operations defined for regular expressions, much as there are for arithmetic expressions and logical operations. Rather than teaching this order to you, we will fully parenthesize our expressions so there is no ambiguity about the order of operations.